# The Computation of Elementary Unitary Matrices 

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#### Abstract

The construction of elementary unitary matrices that transform a complex vector to a real multiple of $e_{1}$, the first column of the identity matrix, are studied. We survey the two well known forms and present what appears to be relatively unknown third form implemented by LAPACK subroutine CLARFG.


## 1 Introduction

Define the elementary Hermitian (or unitary) matrix

$$
\begin{equation*}
U=I-2 w w^{H}, \tag{1}
\end{equation*}
$$

where $w^{H} w=1$. Easily verified is that $U$ is both Hermitian and unitary and is also refered to as a Householder matrix. The matrix $U$ as defined by (1) is a special case of the more general class of elementary matrices defined by

$$
\begin{equation*}
E(u, v ; \sigma)=I-\sigma u v^{H} . \tag{2}
\end{equation*}
$$

See [2], [3] and [4] for further details. The latter reference is a comprehensive study of elementary matrices, their properties and extensions to block implementations.

In [3] Wilkinson states that it is not possible, in general, to find an elementary matrix $U$ such that

$$
\begin{equation*}
U x=\alpha e_{1}, \tag{3}
\end{equation*}
$$

for a complex vector $x$ with $\alpha$ a real number. Note that $U$ is unitary implies that $|\alpha|=\|x\|_{2}$. If the above relation were satisfied then

$$
x^{H} \alpha e_{1}=x^{H} U x .
$$

The righthand side is a real number since $U$ is Hermitian. But if $e_{1}^{T} x$ has a non-zero imaginary part we have a contradiction. The necessary condition is that $x^{H} e_{1}=e_{1}^{T} x$. Equivalently $\epsilon_{1}^{T} x=\xi_{1}$ must have zero imaginary part.

The purpose of this note is to review and examine the case of constructing elementary unitary matrices that satisfy (3) for a complex vector $x$. The next section

[^0]discusses the approach suggested by Wilkinson. The third section introduces yet another approach due to Hammarling and Du Croz. It has been implemented in the NAG Fortran Library subroutine F06HRF, see [1]. A slightly different implementation is the LAPACK subroutine CLARFG. The details of the software implementation for CLARFG are in section four. In fact what led to this study is attempting to understand the differences between the Wilkinson approach and the alternate formulation implemented by LAPACK, [6]. Section five generalizes the Hammarling-Du Croz form for transforming a complex vector $x$ into another complex vector $y$. The last section contains the source code for subroutine CLARFG.

## 2 The Wilkinson Approach

Wilkinson suggests the following modification in Chapter 2, section 45 of [3]. If $\epsilon_{1}^{T} x=\xi_{1}$ has non-zero imaginary part then

$$
\begin{aligned}
x & =e^{i \theta_{1}}\left[\left|\xi_{1}\right|, e^{-i \theta_{1}} \xi_{2}, \ldots, e^{-i \theta_{1}} \xi_{n}\right]^{T}, \\
& =e^{i \theta_{1}} y,
\end{aligned}
$$

where $e^{i \theta_{1}}\left|\xi_{1}\right|$ is the polar form for $\xi_{1}$. Since $\epsilon_{1}^{T} y$ is a real number, a Householder matrix $P$ may be constructed satisfying (3). Set $U=e^{-i \theta_{1}} P$ and

$$
\begin{aligned}
U x & =e^{-i \theta_{1}} P x, \\
& =e^{-i \theta_{1}} P e^{i \theta_{1}} y, \\
& =P y, \\
& =\alpha \epsilon_{1} .
\end{aligned}
$$

The matrix $U$ is not Hermitian but is unitary which is the crucial property. The value of $\alpha$ is not affected;

$$
\begin{aligned}
|\alpha| & =\|U x\|_{2} \\
& =\left|e^{-i \theta_{1}}\right|\|P x\|_{2} \\
& =\|x\|_{2}
\end{aligned}
$$

The matrix $U$ constructed is not an elementary matrix. If $\xi_{1}$ is real then $\theta_{1}$ is arbitrary and may be chosen equal to zero. The resulting elementary matrix is then also Hermitian.

The Wilkinson approach is used by LINPACK when working with complex data.

## 3 An Alternate Approach

The final form for an elementary unitary matrix studied is due to Hammarling and Du Croz, see [1] (Introduction - F06). We now derive this alternate form. Consider the general form (2) for an elementary unitary matrix $U=E(u, v ; \sigma)$. The matrix must be unitary from which it follows that

$$
\begin{aligned}
I & =U^{H} U, \\
& =\left(I-\sigma u v^{H}\right)^{H}\left(I-\sigma u v^{H}\right), \\
& =I-\bar{\sigma} v u^{H}-\sigma u v^{H}+\sigma \bar{\sigma}\left(u^{H} u\right) v v^{H},
\end{aligned}
$$

and cancelling terms results in

$$
\begin{equation*}
\sigma \bar{\sigma}\left(u^{H} u\right) v v^{H}=\bar{\sigma} v u^{H}+\sigma u v^{H} . \tag{4}
\end{equation*}
$$

Rearranging terms gives

$$
\left(\sigma \bar{\sigma}\left(u^{H} u\right) v-\sigma u\right) v^{H}=\bar{\sigma} v u^{H},
$$

and a row space argument implies that $u$ and $v$ are linearly dependent. Substituting $u=v$ into (4) results in

$$
\begin{align*}
|\sigma|^{2}\|v\|_{2}^{2} & =\sigma+\bar{\sigma}  \tag{5}\\
& =2 \operatorname{Re}(\sigma)
\end{align*}
$$

determining the required relationship between $\sigma$ and $v$. Choosing $v=x-\alpha e_{1}$ we have

$$
\begin{aligned}
U x & =\left(I-\sigma v v^{H}\right) x, \\
& =x-\left(\sigma v^{H} x\right) v, \\
& =\alpha \epsilon_{1},
\end{aligned}
$$

if $\sigma^{-1}=v^{H} x$. This choice of $\sigma$ will satisfy (5) as we now demonstrate. First

$$
\begin{aligned}
v^{H} x & =\left(x^{H}-\alpha e_{1}^{T}\right) x, \\
& =x^{H} x-\alpha \xi_{1}, \\
& =\alpha\left(\alpha-\xi_{1}\right),
\end{aligned}
$$

which determines $\sigma$ and

$$
\begin{aligned}
\|v\|_{2}^{2} & =v^{H} v \\
& =\left(x^{H}-\alpha e_{1}^{T}\right)\left(x-\alpha e_{1}\right), \\
& =2 \alpha\left(\alpha-\operatorname{Re}\left(\xi_{1}\right)\right),
\end{aligned}
$$

follows. Finally

$$
\begin{aligned}
\left(v^{H} x\right)\left(x^{H} v\right)(\sigma+\bar{\sigma}) & =\left(v^{H} x\right)\left(x^{H} v\right)\left(\frac{1}{v^{H} x}+\frac{1}{x^{H} v}\right), \\
& =x^{H} v+v^{H} x \\
& =\alpha\left(\alpha-\xi_{1}\right)+\alpha\left(\alpha-\bar{\xi}_{1}\right) \\
& =2 \alpha\left(\alpha-\operatorname{Re}\left(\xi_{1}\right)\right)
\end{aligned}
$$

shows

$$
\frac{\sigma+\bar{\sigma}}{|\sigma|^{2}}=\|v\|_{2}^{2}
$$

as claimed. The form as outlined in this section does not appear to be as widely known as the Wilkinson one. It has the benefit of transforming a complex vector $x$ directly to real multiple of $e_{1}$. When $\xi_{1}$ is purely real then $U$ is Hermitian. Unlike the Wilkinson variant the Hammarling-Du Croz approach results in an elementary unitary matrix.

## 4 Software Implementation of CLARFG

The actual implementation of the Hammarling-Du Croz variant has some slight modifications. The resulting code is an excellent example of the art of developing software from a numerical algorithm. Subroutine CLARFG determines an elementary unitary matrix $U=I-\tau u u^{H}$ such that $U^{H} x=\alpha e_{1}$ where $|\alpha|=\|x\|_{2}$ for $\alpha$ a real number. Simplifying the implementation of other algorithms in LAPACK requiring the use of elementary unitary matrices the normalization $u^{H} e_{1}=1$ is taken. This normalization is done for storage considerations and is discussed in [5]. From the previous section it follows that

$$
\begin{aligned}
U^{H} & =I-\frac{1}{\alpha\left(\alpha-\bar{\xi}_{1}\right)} v v^{H}, \\
& =I-\frac{\left(\xi_{1}-\alpha\right)\left(\bar{\xi}_{1}-\alpha\right)}{\alpha\left(\alpha-\bar{\xi}_{1}\right)} u u^{H}, \\
& =I-\frac{\alpha-\xi_{1}}{\alpha} u u^{H}, \\
& =I-\tau u u^{H},
\end{aligned}
$$

where

$$
\begin{aligned}
u & =\frac{v}{\xi_{1}-\alpha} \\
& =\frac{x-\alpha \epsilon_{1}}{\xi_{1}-\alpha}
\end{aligned}
$$

and

$$
\tau=\frac{\alpha-\xi_{1}}{\alpha}
$$

Storage of $U$ for use in further computation only requires storage for the complex $\tau$ and $x$ may be overwritten with both $\alpha$ and the essential part of $u$, i.e.

$$
x \leftarrow\left[\alpha, \frac{\xi_{2}}{\xi_{1}-\alpha}, \ldots, \frac{\xi_{n}}{\xi_{1}-\alpha}\right]^{T} .
$$

In order that the value for $\tau$ and the scaling factor $\xi_{1}-\alpha$ have small relative error

$$
\alpha \leftarrow-\operatorname{sign}\left(\operatorname{Re}\left(\xi_{1}\right)\right)\|x\|_{2},
$$

is chosen. This is a standard modification and is mentioned for the sake of completeness. The resulting $\tau$ satisfies the two properties

$$
\begin{array}{rc}
1 & \leq \operatorname{Re}(\tau) \leq 2, \\
|\tau-1| & \leq \\
1,
\end{array}
$$

when $x \neq \gamma e_{1}$ with $\gamma$ real. If $x$ is a real multiple of $e_{1}$ then set

$$
\begin{aligned}
\tau & \leftarrow 0 \\
U & \leftarrow I
\end{aligned}
$$

The reviewer of CLARFG will notice one final point which needs explanation. The careful programmer took care not to reciprocate the number $\|x\|_{2}$ that may fall below a certain machine dependent tolerance, SAFMIN. The value SAFMIN, computed by the LAPACK auxiliary subroutine SLAMCH is a machine dependent lower bound for numbers that may be safely reciprocated and not cause an overflow condition. If $\|x\|_{2}$ is less than the lower bound then the vector $x$ is scaled by a multiple of the reciprocal of SAFMIN until $\|\theta x\|_{2}$ is at least as large as SAFMIN. Defining the integer $k$ to represent the number of scalings required results in

$$
\theta=k \frac{1}{\text { SAFMIN }}
$$

The number $\tau$ may now be safely computed as

$$
\tau \leftarrow \frac{\|\theta x\|_{2}-\theta \xi_{1}}{\|\theta x\|_{2}}
$$

In a similar fashion

$$
\alpha \leftarrow-\operatorname{sign}\left(\operatorname{Re}\left(\xi_{1}\right)\right) \frac{1}{\theta}\left(\|\theta x\|_{2}\right)
$$

and the essential part of $u$

$$
u \leftarrow \frac{1}{\theta \xi_{1}-\theta \alpha}\left[\theta \xi_{2}, \ldots, \theta \xi_{n}\right]^{T}
$$

are computed.

## 5 Generalizations

We conclude with the problem of determining an elementary unitary matrix $U$ that satisfies

$$
\begin{align*}
U x & =\left(I-\sigma v v^{H}\right) x  \tag{6}\\
& =y
\end{align*}
$$

where $\|x\|_{2}=\|y\|_{2}$. The derivation of section three requires little modification. The value of $\sigma^{-1}=v^{H} x$ and

$$
\begin{aligned}
v & =x-y \\
v^{H} x & =\|x\|_{2}^{2}-y^{H} x
\end{aligned}
$$

may be shown to satisfy (5) and (6).

## 6 Acknowledgements

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## References

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[2] A. H. Householder, The Theory of Matrices in Numerical Analysis, Blaisdell, 1964.
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[4] A. A. Dubrulle, Work Notes on Elementary Matrices, Hewlett-Packard Laboratories, 1993, Technical Report HPL-93-69.
[5] G. H. Golub, C. F. Van Loan, Matrix Computations, Johns Hopkins, second edition, 1989.
[6] E. Anderson, Z. Bai, C. Bischoff, J. Demmel, J. Dongarra, J. Du Croz, A. Greenbaum, S. Hammarling, A. McKenney, S. Ostrouchov, and D. C. Sorensen, LAPACK Users' Guide, SIAM, 1992.
[7] J. J. Dongarra, C. B. Moler, J. R. Bunch, G. W. Stewart, LINPACK Users' Guide, SIAM, 1979.

## 7 LAPACK Subroutine CLARFG

```
    SUBROUTINE CLARFG( N, ALPHA, X, INCX, TAU )
```

* 
* -- LAPACK auxiliary routine (version 1.1) --
* Univ. of Tennessee, Univ. of California Berkeley, NAG Ltd.,
* Courant Institute, Argonne National Lab, and Rice University
* February 29, 1992
* 
* .. Scalar Arguments ..
INTEGER INCX, N
COMPLEX ALPHA, TAU
* . .
* .. Array Arguments . .
COMPLEX X ( * )
. .
* 
* Purpose
* =======
* 
* CLARFG generates a complex elementary reflector $H$ of order $n$, such
* that
* 
* $H^{\prime} *($ alpha $)=($ beta $), H^{\prime} * H=I$.
* ( x$) \quad(0)$
* 

```
where alpha and beta are scalars, with beta real, and x is an
(n-1)-element complex vector. H is represented in the form
    H=I - tau * ( 1 ) * ( 1 v' ) ,
    ( v )
where tau is a complex scalar and v is a complex (n-1)-element
vector. Note that H is not hermitian.
If the elements of x are all zero and alpha is real, then tau = 0
and H is taken to be the unit matrix.
Otherwise 1 <= real(tau) <= 2 and abs(tau-1) <= 1 .
Arguments
=========
N (input) INTEGER
    The order of the elementary reflector.
    (input/output) COMPLEX
    On entry, the value alpha.
    On exit, it is overwritten with the value beta.
    (input/output) COMPLEX array, dimension
                                    (1+(N-2)*abs(INCX))
    On entry, the vector x.
    On exit, it is overwritten with the vector v.
    (input) INTEGER
    The increment between elements of X. INCX <> 0.
        (output) COMPLEX
        The value tau.
```



```
    .. Parameters ..
    REAL ONE, ZERO
    PARAMETER ( ONE = 1.OE+0, ZERO = 0.OE+0)
    .. Local Scalars ..
    INTEGER J, KNT
    REAL ALPHI, ALPHR, BETA, RSAFMN, SAFMIN, XNORM
    . External Functions ..
    REAL
                            SCNRM2, SLAMCH, SLAPY3
```

```
    COMPLEX CLADIV
    EXTERNAL CLADIV, SCNRM2, SLAMCH, SLAPY3
* ..
    .. Intrinsic Functions ..
    INTRINSIC ABS, AIMAG, CMPLX, REAL, SIGN
* .. External Subroutines ..
    EXTERNAL CSCAL, CSSCAL
    .. Executable Statements ..
    IF( N.LE.O ) THEN
        TAU = ZERO
        RETURN
    END IF
*
    XNORM = SCNRM2( N-1, X, INCX )
    ALPHR = REAL( ALPHA )
    ALPHI = AIMAG( ALPHA )
    IF( XNORM.EQ.ZERO .AND. ALPHI.EQ.ZERO ) THEN
        H = I
            TAU = ZERO
        ELSE
    general case
        BETA = -SIGN( SLAPY3( ALPHR, ALPHI, XNORM ), ALPHR )
        SAFMIN = SLAMCH( 'S' )
        RSAFMN = ONE / SAFMIN
            IF( ABS( BETA ).LT.SAFMIN ) THEN
        XNORM, BETA may be inaccurate; scale X and recompute them
        KNT = 0
    10 CONTINUE
        KNT = KNT + 1
        CALL CSSCAL( N-1, RSAFMN, X, INCX )
        BETA = BETA*RSAFMN
        ALPHI = ALPHI*RSAFMN
        ALPHR = ALPHR*RSAFMN
        IF( ABS( BETA ).LT.SAFMIN )
$ GO TO 10
```

```
*
*
    XNORM = SCNRM2( N-1, X, INCX )
    ALPHA = CMPLX( ALPHR, ALPHI )
    BETA = -SIGN( SLAPY3( ALPHR, ALPHI, XNORM ), ALPHR )
    TAU = CMPLX( ( BETA-ALPHR ) / BETA, -ALPHI / BETA )
    ALPHA = CLADIV( CMPLX( ONE ), ALPHA-BETA )
    CALL CSCAL( N-1, ALPHA, X, INCX )
    If ALPHA is subnormal, it may lose relative accuracy
    ALPHA = BETA
    DO 20 J = 1, KNT
        ALPHA = ALPHA*SAFMIN
    CONTINUE
        ELSE
    TAU = CMPLX( ( BETA-ALPHR ) / BETA, -ALPHI / BETA )
    ALPHA = CLADIV( CMPLX( ONE ), ALPHA-BETA )
    CALL CSCAL( N-1, ALPHA, X, INCX )
    ALPHA = BETA
        END IF
END IF
RETURN
End of CLARFG
END
```


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